

conditional and marginal (individual), informations associated with the bivariate probability distribution ( $p_{ij}$ ).

We may define the marginal probability distributions of  $X$  and  $Y$  by

$$p_{i0} = \sum_{j=1}^n p_{ij} \text{ and } p_{0j} = \sum_{i=1}^m p_{ij} \text{ for all } i, j.$$

Then, obviously the marginal entropies of the two marginal distributions are given by

$$H(X) = - \sum_{i=1}^m p_{i0} \log p_{i0} \text{ and } H(Y) = - \sum_{j=1}^n p_{0j} \log p_{0j}.$$

The entropy  $H(X)$  measures the uncertainty of the message sent (irrespective of the message received) and  $H(Y)$  performs the same role for the message received.

The joint entropy is the entropy of the joint distribution of the messages sent and received, and is therefore given by

$$H(X, Y) = - \sum_{i=1}^m \sum_{j=1}^n p_{ij} \log p_{ij}.$$

It may be observed that :

$$\text{Max } H(X, Y) = \log mn = \log m + \log n = \text{max } H(X) + \text{max } H(Y).$$

**Theorem 26.2.**  $H(X, Y) \leq H(X) + H(Y)$  with equality, if and only if,  $X$  and  $Y$  are independent.

**Proof.** We may write

$$\begin{aligned} H(X) + H(Y) &= - \sum_{i=1}^m p_{i0} \log p_{i0} - \sum_{j=1}^n p_{0j} \log p_{0j} \\ &= - \sum_{i=1}^m \left( \sum_{j=1}^n p_{ij} \right) \log p_{i0} - \sum_{j=1}^n \left( \sum_{i=1}^m p_{ij} \right) \log p_{0j} \\ &= - \sum_{i=1}^m \sum_{j=1}^n p_{ij} \log (p_{i0} p_{0j}) = - \sum_{i=1}^m \sum_{j=1}^n p_{ij} \log q_{ij}, \text{ where } q_{ij} = p_{i0} p_{0j}. \end{aligned} \quad \dots(i)$$

Also, by definition 
$$H(X, Y) = - \sum_{i=1}^m \sum_{j=1}^n p_{ij} \log p_{ij}. \quad \dots(ii)$$

But, we observe that

$$\sum_{i=1}^m \sum_{j=1}^n q_{ij} = \sum_{i=1}^m \sum_{j=1}^n p_{i0} p_{0j} = \left( \sum_{i=1}^m p_{i0} \right) \left( \sum_{j=1}^n p_{0j} \right) = 1 = \sum_{i=1}^m \sum_{j=1}^n p_{ij}$$

By virtue of Theorem 26.1, (i) and (ii), it follows that

$$H(X, Y) \leq H(X) + H(Y)$$

with equality, if and only if,  $q_{ij} = p_{ij}$  for all  $i$  and  $j$ . The condition for equality reduces to  $p_{i0} p_{0j} = p_{ij}$  meaning thereby  $X$  and  $Y$  are independent.

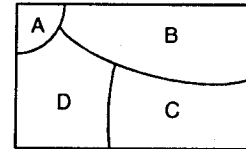


Fig. 26.6

- Q. 1.** Show that  $H(X)$  achieves its maximum if all the values of  $X$  are equi-probable. [Delhi (OR) 92]  
**2.** Let  $X$  and  $Y$  be two discrete random variables, each taking a finite number of values. Show that  $H(Y|X) \leq H(Y)$  with equality iff  $X$  and  $Y$  are independent

**26.10-2. Conditional Entropies**

Consider two finite discrete sample spaces  $S_1$  and  $S_2$  and their product space  $S$ . In  $S_1$  and  $S_2$ , select complete sets of events in the sense of equations (26.17) and (26.18).

$$\{E\} = [E_1, E_2, \dots, E_n] \quad \dots(26.26)$$

$$\{F\} = [F_1, F_2, \dots, F_m] \quad \dots(26.27)$$

Thus, the complete set of events in the product space  $S_1 \times S_2$  will be given by

$$\{EF\} = \begin{bmatrix} E_1F_1 & E_1F_2 & E_1F_3 & \dots & E_1F_m \\ E_2F_1 & E_2F_2 & E_2F_3 & \dots & E_2F_m \\ \vdots & \vdots & \vdots & \dots & \vdots \\ E_nF_1 & E_nF_2 & E_nF_3 & \dots & E_nF_m \end{bmatrix} \quad \dots(26.28)$$

[By taking cartesian product of two sets  $\{E\}$  and  $\{F\}$ ]

For example, an event  $F_r$  may occur in conjunction with  $E_1, E_2, \dots, \text{ or } E_n$ .

$$F_r = \sum_{k=1}^n E_k E_r \quad \dots(26.29)$$

$$P\{X = x_k | Y = y_r\} = \frac{P\{X = x_k \cap Y = y_r\}}{P\{Y = y_r\}} \quad \dots(26.30)$$

or

$$p\{x_k | y_r\} = \frac{p\{k, r\}}{p\{y_r\}} \quad \dots(26.31)$$

Now, consider the following probability scheme

$$\{E | F_r\} = [E_1 | E_r, E_2 | E_r, E_3 | E_r, \dots, E_n | E_r] \quad \dots(26.32)$$

$$P\{E | F_r\} = \left[ \frac{p\{1, r\}}{p\{y_r\}}, \frac{p\{2, r\}}{p\{y_r\}}, \frac{p\{3, r\}}{p\{y_r\}}, \dots, \frac{p\{n, r\}}{p\{y_r\}} \right] \quad \dots(26.33)$$

The probability scheme is not only finite but also complete because the sum of elements of this matrix is unity. Therefore, an entropy will be given by

$$H\{X | y_r\} = - \sum_{k=1}^n \frac{p\{k, r\}}{p\{y_r\}} \log \frac{p\{k, r\}}{p\{y_r\}} \quad \dots(26.34)$$

$$= - \sum_{k=1}^n p\{x_k | y_r\} \log p\{x_k | y_r\} \quad \dots(26.35)$$

Now, take the average of this conditional entropy for all admissible values of  $y_r$ , so that a measure of average conditional entropy of the system can be obtained.

$$\begin{aligned} H\{X | Y\} &= \overline{H\{X | y_r\}} = \sum_{r=1}^m p\{y_r\} \cdot H\{X | y_r\} \\ &= - \sum_{r=1}^m p\{y_r\} \sum_{k=1}^n p\{x_k | y_r\} \log p\{x_k | y_r\} \end{aligned} \quad \dots(26.36)$$

$$H\{X | Y\} = - \sum_{r=1}^m \sum_{k=1}^n p\{y_r\} \cdot p\{x_k | y_r\} \log p\{x_k | y_r\} \quad \dots(26.37)$$

Likewise, it is possible to obtain the expression for the average conditional entropy  $H\{Y | X\}$ , i.e.

$$H\{Y | X\} = - \sum_{k=1}^n \sum_{r=1}^m p\{x_k\} \cdot p\{y_r | x_k\} \log p\{y_r | x_k\} \quad \dots(26.38)$$

Two conditional entropies may be expressed as

$$H\{X | Y\} = - \sum_{r=1}^m \sum_{k=1}^n p\{x_k, y_r\} \log p\{x_k | y_r\} \quad \dots(26.39)$$

$$H\{Y | X\} = - \sum_{k=1}^n \sum_{r=1}^m p\{x_k, y_r\} \log p\{y_r | x_k\} \quad \dots(26.40)$$

Note. It should be noted that all entropies are positive numbers as they are sum of positive numbers.

**26.11. AN IMPORTANT THEOREM**

**Theorem 26.3.** Prove that

$$H(X, Y) = H(X|Y) + H(Y) = H(Y|X) + H(X),$$

where  $H(X) \geq H(X|Y)$ .

[Delhi (O.R.) 92]

**Proof.** By definition of entropy function  $H$ ,

$$H(X) = - \sum_{k=1}^n p\{x_k\} \log p\{x_k\} \quad \dots(26.41)$$

$$H(Y) = - \sum_{r=1}^m p\{y_r\} \log p\{y_r\} \quad \dots(26.42)$$

The joint entropy function of  $X$  and  $Y$  will be given by

$$H(X, Y) = - \sum_{k=1}^n \sum_{r=1}^m p\{k, r\} \log p\{k, r\} \quad \dots(26.43)$$

where  $p(k, r)$  is the joint probability for the occurrence of two events  $E_k$  and  $E_r$  simultaneously.

Rewriting expressions (26.39) and (26.40) for conditional probabilities, we have

$$H(X|Y) = - \sum_{r=1}^m \sum_{k=1}^n p\{x_k, y_r\} \log p\{x_k|y_r\} \quad \dots(26.39a)$$

$$H(Y|X) = - \sum_{k=1}^n \sum_{r=1}^m p\{x_k, y_r\} \log p\{y_r|x_k\} \quad \dots(26.40a)$$

Now, basic relationships among *marginal, joint and conditional* probabilities are :

$$p\{x_k, y_r\} = p\{x_k|y_r\} p\{y_r\} = p\{y_r|x_k\} p\{x_k\} \quad \dots(26.44)$$

$$\log p\{x_k, y_r\} = \log p\{x_k|y_r\} + \log p\{y_r\} = \log p\{y_r|x_k\} + \log p\{x_k\} \quad \dots(26.45)$$

(i) To prove  $H(X, Y) = H(Y|X) + H(X)$  and  $H(X, Y) = H(X|Y) + H(Y)$ .

Directly substituting the relations (26.44) and (26.45) in defining equations (26.41), (26.42), (26.39a), (26.43) and (26.40a) of entropies, these two basic identities can be easily proved.

(ii) To prove  $H(X) \geq H(X|Y)$ .

For the proof of this inequality,

$$\begin{aligned} H(X|Y) - H(X) &= \sum_{r=1}^m \sum_{k=1}^n p\{x_k, y_r\} \log \frac{p\{x_k\}}{p\{x_k|y_r\}} \\ &\leq \sum_{r=1}^m \sum_{k=1}^n p\{x_k, y_r\} \left[ \frac{p\{x_k\}}{p\{x_k|y_r\}} - 1 \right] \log e \quad \dots(26.46) \end{aligned}$$

(since  $\log x \leq (x - 1) \log e$  by convexity of logarithmic function).

But, 
$$\sum_{r=1}^m \sum_{k=1}^n [p\{x_k\} \cdot p\{y_r\} - p\{x_k, y_r\}] \log e = \sum_{r=1}^m [p\{y_r\} - p\{y_r\}] \log e = 0$$

Hence,  $H(X|Y) - H(X) \leq 0$  or  $H(X) \geq H(X|Y)$ .

Similarly, it can be proved that  $H(Y) \geq H(Y|X)$ .

**Q. 1.** Define conditional entropy. In the usual notation show that  $H(Y) + H(X) \geq H(X, Y)$ .

**2.** Show that average amount of information increases if one of the events is partitioned.

[Delhi (OR) 92]

**3.** Let  $X$  and  $Y$  be two discrete random variables, each taking finite number of values. Prove that

$$H(X, Y) = H(X) + H(Y|X) = H(Y) + H(X|Y)$$

[Delhi (OR) 92]

**Example 10.** A transmitter has an alphabet consisting of five letters  $\{x_1, x_2, x_3, x_4, x_5\}$  and the receiver has an alphabet consisting of four letters  $\{y_1, y_2, y_3, y_4\}$ . The joint probabilities for the communication are given below :

	$y_1$	$y_2$	$y_3$	$y_4$		
$x_1$	[	0.25	0.00	0.00	0.00	]
$x_2$	[	0.10	0.30	0.00	0.00	]
$x_3$	[	0.00	0.05	0.10	0.00	]
$x_4$	[	0.00	0.00	0.05	0.10	]
$x_5$	[	0.00	0.00	0.05	0.00	]

Determine the Marginal, Conditional and Joint entropies for this channel (Assume  $0 \log 0 \equiv 0$ ).

**Solution.** The channel is described here by the joint probabilities  $p_{ij}$ ,  $i = 1, 2, \dots, 5$  and  $j = 1, 2, \dots, 4$ . The conditional and marginal probabilities are easily obtained from  $p_{ij}$ 's as given below :

$$p_{10} = 0.25 + 0.00 = 0.25, \quad p_{20} = 0.10 + 0.30 = 0.40, \quad p_{30} = 0.05 + 0.10 = 0.15,$$

$$p_{40} = 0.05 + 0.10 = 0.15, \quad p_{50} = 0.05 + 0.00 = 0.05.$$

Similarly,  $p_{01} = 0.35$ ,  $p_{02} = 0.35$ ,  $p_{03} = 0.20$ ,  $p_{04} = 0.10$ .

By using the result ( $p_{j|i} = p_{ij}/p_{i0}$ ), the conditional probabilities are given in the following channel matrix :

**Conditional Prob. Matrix ( $p_{j|i}$ )**

	1	2	3	4		
1	[	1	0	0	0	]
2	[	1/4	3/4	0	0	]
3	[	0	1/3	2/3	0	]
4	[	0	0	1/3	2/3	]
5	[	0	0	1	0	]

Now various entropies associated with this channel are obtained (taking all logarithms to the base 2).

**Marginal Entropies :**

$$\begin{aligned} H(X) &= - \sum_{i=1}^5 p_{i0} \log p_{i0} \\ &= - [(.25) \log (.25) + (.40) \log (.40) + 2 (.15) \log (.15) + (.05) \log (.05)] \\ &= \frac{1}{4} \log 4 + \frac{2}{5} \log \frac{5}{2} + \frac{3}{10} \log \frac{20}{3} + \frac{1}{20} \log 20 = 1.3260 \text{ bits.} \end{aligned}$$

$$\begin{aligned} H(Y) &= - \sum_{j=1}^4 p_{0j} \log p_{0j} = - [2 (.35) \log (.35) + (.20) \log (.20) + (.10) \log (.10)] \\ &= \frac{7}{10} \log (2.857) + \frac{1}{5} \log 5 + \frac{1}{10} \log 10 \\ &= 1.8556 \text{ bits} \end{aligned}$$

**Conditional Entropies :**

$$\begin{aligned} H(Y|X) &= - \sum_{i=1}^5 \sum_{j=1}^4 p_{ij} \log p_{j|i} \\ &= (.25) \log 1 + (.10) \log 4 + (.30) \log (4/3) + (.05) \log 3 \\ &\quad + (.10) \log (3/2) + (.05) \log 3 + (.10) \log (3/2) + (.05) \log 1 \\ &= \frac{1}{10} \log 4 + \frac{3}{10} \log 4 - \frac{3}{10} \log 3 + \frac{1}{10} \log 3 + \frac{1}{5} \log 3 - \frac{1}{5} \log 2 \\ &= \frac{1}{10} [4 \log 4 - 2 \log 2] = 6/10 = 0.60 \text{ bits.} \end{aligned}$$

By Theorem 26.3, we have

$$H(X|Y) = H(X) + H(Y|X) - H(Y) = 1.3260 + 0.6000 - 1.8556 = 0.0704 \text{ bits.}$$

**Joint Entropies :**

$$H(X, Y) = H(X) + H(Y|X) = 1.3260 + 0.6000 = 1.9260 \text{ bits.}$$

### 26.12. SET OF AXIOMS FOR AN ENTROPY FUNCTION

Assume the following four conditions as axioms :